# Completeness and synchronization for finite sets of words 

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#### Abstract

We investigate the problem of finding upper bounds to the minimal lengths of incompletable and synchronizing words of a finite language $X$ in terms of the length of the words of $X$. This problem is related to two well-known conjectures formulated by Černý and Restivo respectively.


Keywords: Černý conjecture, synchronizing automaton, incompletable word, synchronizing set, complete set

The concepts of completeness and synchronization play a central role in Computer Science since they appear in the study of several problems on variable length codes and on finite automata.
We recall that a set of words $X$ is incomplete if there exists a word which is not factor of any word of $X^{*}$. Such a word is called incompletable. The set $X$ is synchronizing if there is a pair of words $u, v \in X^{*}$ such that whenever ruvs $\in X^{*}$, one has ru, vs $\in X^{*}$. Such a pair is called synchronizing.
According to a well-known result of Schützenberger, the property of completeness provides an algebraic characterization of finite maximal codes, which are the objects used in Information Theory to construct optimal sequential codings. A problem that naturally arises in this context concerns the minimal length of an incompletable word in a finite incomplete set of words. This problem was first addressed and studied by Restivo in [20]. In particular, it was conjectured that a finite incomplete set $X$ has always an incompletable word whose length is quadratic in the size of $X$, that is, the maximal length of the words of $X$. Results on this problem have been obtained in $[5,15,16,20]$.
The property of synchronization plays a natural role for the formal description of fault-tolerant systems. In particular, in the context of Information Theory, this property is relevant for the construction of decoders that are able to efficiently cope with decoding errors caused by noise during the data transmission. In the study of synchronizing sets, the search for synchronizing words of minimal length in a prefix complete code is tightly related to that of minimal length synchronizing words of synchronizing deterministic automata and the celebrated Černý Conjecture [13] (see also $[1,2,3,4,6,7,9,10,11,12,13,14,17,18,19,21]$ for some results on the problem). In particular, in [2] (see also [3]), Béal and Perrin have proved that a complete synchronizing prefix code $X$ on an alphabet of $d$ letters with $n$ code-words has a synchronizing word of length $O\left(n^{2}\right)$.

[^0]In this paper we are interested in finding upper bounds to the minimal lengths of incompletable and synchronizing words of a finite set $X$ in terms of the size of $X$.
Let $\mathcal{L}$ be a class of finite languages and denote by $\ell(X)$ the size of a finite set $X$ of words. For all $n, d>0$, we denote by $R_{\mathcal{L}}(n, d)$ the least positive integer $r$ satisfying the following condition: any incomplete set $X \in \mathcal{L}$ on a $d$-letter alphabet such that $\ell(X) \leq n$ has an incompletable word of length $r$.

Similarly, we denote by $C_{\mathcal{L}}(n, d)$ the least positive integer $c$ satisfying the following condition: any synchronizing set $X \in \mathcal{L}$ on a $d$-letter alphabet such that $\ell(X) \leq n$ has a synchronizing pair $(u, v)$ such that $|u v| \leq c$.
In this context, our main result provides a bridge between the parameters $R_{\mathcal{L}}(n, d)$ and $C_{\mathcal{L}}(n, d)$.

Proposition 1 Let $\mathcal{F}$ and $\mathcal{M}$ be, respectively, the classes of finite languages and of complete finite codes. For all $n, d>0$ one has

$$
C_{\mathcal{M}}(n, d) \leq 2 R_{\mathcal{F}}(n, d+1)+2 n-2 .
$$

In particular, if Restivo's conjecture is true, Proposition 1 gives

$$
C_{\mathcal{M}}(n, d)=O\left(n^{2}\right),
$$

thus providing a quadratic bound in the size of the set for the minimal length of a synchronizing pair of a finite synchronizing complete code.
We notice that, in rather general cases, quadratic lower bounds for $R_{\mathcal{L}}$ and $C_{\mathcal{L}}$ are known. Indeed, if $\mathcal{P}$ is the class of prefix codes, then

$$
R_{\mathcal{P}}(n, 2) \geq n^{2}+n-1, \quad C_{\mathcal{P}}(n, 2) \geq(n-1)^{2} .
$$

The first inequality is proved in [19], the second follows from a result of [1].
We outline the proof of Proposition 1 in the particular case that $X$ is a prefix code. Let $\mathcal{A}$ be the minimal deterministic finite automaton accepting $X^{*}, q_{0}$ be its initial state, and $a$ be the terminal letter of some word of $X$. We construct a non deterministic finite automaton $\mathcal{A}^{\prime}$ adding to $\mathcal{A}$ a new letter $a^{\prime}$ with transitions defined by $q a^{\prime}=\left\{q a, q_{0}\right\}$ if $q a \neq q_{0}, q a^{\prime}=\emptyset$, otherwise. It turns out that the language accepted by $\mathcal{A}^{\prime}$ is equal to $Y^{*}$ for a suitable incomplete set $Y$ of the same size of $X$. Moreover, replacing by $a$ all the occurrences of $a^{\prime}$ in a minimal incompletable word of $Y$, one obtains a synchronizing word of $X$. This allows to establish, for the class $\mathcal{P}$ of complete prefix codes, the inequality

$$
C_{\mathcal{P}}(n, d) \leq R_{\mathcal{F}}(n, d+1) .
$$

For instance, if $X=\{a, b a a b, b a b, b b\}$, then $\mathcal{A}$ and $\mathcal{A}^{\prime}$ are respectively the automata below, a minimal incompletable word of $Y$ is $a a a^{\prime}$, and $a a a$ is a synchronizing word of the prefix code $X$.


The proof of the general case [8] makes use of unambiguous finite state automata and of their combinatorial properties.

Afterwards, we study the dependence of the parameters $R_{\mathcal{L}}(n, d)$ and $C_{\mathcal{L}}(n, d)$ upon the number of letters $d$ of the considered alphabet, by showing that both the parameters have a low rate of growth.

Proposition 2 Let $\mathcal{L}$ be the class of finite languages (resp. codes, prefix codes). Then, for all $d>2$ one has

$$
R_{\mathcal{L}}(n, d) \leq\left\lceil\frac{R_{\mathcal{L}}\left(\left\lceil\log _{2} d\right\rceil n, 2\right)}{\left\lfloor\log _{2} d\right\rfloor}\right\rceil, \quad C_{\mathcal{L}}(n, d) \leq\left\lceil\frac{C_{\mathcal{L}}\left(\left\lceil\log _{2}(d+1)\right\rceil n, 2\right)}{\left\lceil\log _{2}(d+1)\right\rceil}\right\rceil
$$

Proposition 3 Let $\mathcal{L}$ be the class of complete finite languages (resp. codes, prefix codes). Then, for all $d>2$ one has

$$
C_{\mathcal{L}}(n, d) \leq\left\lceil\frac{C_{\mathcal{L}}\left(\left\lceil\log _{2}(d+1)\right\rceil n, 2\right)}{\left\lfloor\log _{2}(d-1)\right\rfloor}\right\rceil
$$

The proofs of the propositions above are based on suitable binary encodings preserving incompleteness or synchronization. Among the combinatorial tools used in this work, we mention in particular the following one.

Lemma 1 Let $d, n \geq 1$. The positive integers $k_{1}, \ldots, k_{n}$ are the code-word lengths of a synchronizing complete prefix code over $d$ letters if and only if

$$
\operatorname{gcd}\left(k_{1}, k_{2}, \ldots, k_{n}\right)=1, \quad \sum_{i=1}^{n} d^{-k_{i}}=1
$$

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